

A THEORY OF 2-PRO-OBJECTS

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Résumé. Dans [1] Grothendieck développe la théorie des pro-objets sur une catégorie \mathcal{C} . La propriété fondamentale de la catégorie $\text{Pro}(\mathcal{C})$ des pro-objets est qu'il y a une immersion $\mathcal{C} \xrightarrow{c} \text{Pro}(\mathcal{C})$, $\text{Pro}(\mathcal{C})$ est fermée par petites limites cofiltrées, et ces limites sont libres dans le sens que pour une catégorie quelconque \mathbf{E} fermée par petites limites cofiltrées, la précomposition par c détermine une équivalence des catégories $\text{Cat}(\text{Pro}(\mathcal{C}), \mathbf{E})_+ \simeq \text{Cat}(\mathcal{C}, \mathbf{E})$ (où "+" indique la sous-catégorie des foncteurs qui préservent les limites cofiltrées). Dans cet article nous développons une théorie des pro-objets en dimension 2. Étant donnée une 2-catégorie \mathcal{C} , nous construisons une 2-catégorie $2\text{-Pro}(\mathcal{C})$, dont nous appelons les objets 2-pro-objets. Nous montrons que $2\text{-Pro}(\mathcal{C})$ a toutes les propriétés basiques attendues, correctement relativisées au contexte 2-catégorique, y compris la propriété universelle analogue à celle mentionnée ci-dessus. Bien que nous ayons à notre disposition les résultats de la théorie des catégories enrichies, notre théorie va au-delà du cas des catégories enrichies sur Cat , car nous considérons la notion non-strict de pseudo-limite, qui est usuellement celle d'intérêt pratique.

Abstract. In [1], Grothendieck develops the theory of pro-objects over a category \mathcal{C} . The fundamental property of the category $\text{Pro}(\mathcal{C})$ is that there is an embedding $\mathcal{C} \xrightarrow{c} \text{Pro}(\mathcal{C})$, the category $\text{Pro}(\mathcal{C})$ is closed under small cofiltered limits, and these limits are free in the sense that for any category \mathbf{E} closed under small cofiltered limits, pre-composition with c determines an equivalence of categories $\text{Cat}(\text{Pro}(\mathcal{C}), \mathbf{E})_+ \simeq \text{Cat}(\mathcal{C}, \mathbf{E})$, (where the "+" indicates the full subcategory of the functors preserving cofiltered limits). In this paper we develop a 2-dimensional theory of pro-objects. Given a 2-category \mathcal{C} , we define the 2-category $2\text{-Pro}(\mathcal{C})$ whose

objects we call 2-pro-objects. We prove that $2\text{-Pro}(\mathcal{C})$ has all the expected basic properties adequately relativized to the 2-categorical setting, including the universal property corresponding to the one described above. We have at hand the results of *Cat*-enriched category theory, but our theory goes beyond the *Cat*-enriched case since we consider the non strict notion of pseudo-limit, which is usually that of practical interest.

Key words. 2-pro-object, 2-filtered, pseudo-limit.

MS classification. Primary 18D05, Secondary 18A30.

Introduction. In this paper we develop a 2-dimensional theory of pro-objects. Our motivation are intended applications in homotopy, in particular strong shape theory. The Čech nerve before passing modulo homotopy determines a 2-pro-object which is not a pro-object, leaving outside the actual theory of pro-objects. Also, the theory of 2-pro-objects reveals itself a very interest subject in its own right.

Given a 2-category \mathcal{C} , we define the 2-category $2\text{-Pro}(\mathcal{C})$, whose objects we call 2-pro-objects. A 2-pro-object is a 2-functor (or diagram) indexed in a 2-cofiltered 2-category. Our theory goes beyond enriched category theory because in the definition of morphisms, instead of strict 2-limits, we use the non strict notion of pseudo-limit, which is usually that of practical interest. We prove that $2\text{-Pro}(\mathcal{C})$ has all the expected basic properties of the category of pro-objects, adequately relativized to the 2-categorical setting.

Section 1 contains some background material on 2-categories. Most of this is standard, but some results (for which we provide proofs) do not appear to be in the literature. In particular we prove that pseudolimits are computed pointwise in the 2-functor 2-categories $\mathcal{H}om(\mathcal{C}, \mathcal{D})$ and $\mathcal{H}om_p(\mathcal{C}, \mathcal{D})$ (definition 1.1.11), with 2-natural or pseudonatural transformations as arrows. This result, although expected, needs nevertheless a proof. We recall from [8] the construction of 2-filtered pseudocolimits of categories which is essential for the computations in the 2-category of 2-pro-objects introduced in section 2. Finally, we consider the notion of flexible functors from [4] and state a useful characterization independent of the left adjoint to the inclusion $\mathcal{H}om(\mathcal{C}, \mathcal{D}) \rightarrow \mathcal{H}om_p(\mathcal{C}, \mathcal{D})$ (Proposition 1.3.2). With this characterization the pseudo Yoneda lemma just says

that the representable 2-functors are flexible. It follows also that the 2-functor associated to any 2-pro-object is flexible, and this has important consequences for a Quillen model structure in the 2-category of 2-pro-objects currently being developed by the authors in ongoing research.

Section 2 contains the main results of this paper. In a first subsection we define the 2-category of 2-pro-objects of a 2-category \mathcal{C} and establish the basic formula for the morphisms and 2-cells between 2-pro-objects in terms of pseudo limits and pseudo colimits of the hom categories of \mathcal{C} . With this, inspired in the notion of an arrow representing a morphism of pro-objects found in [3], in the next subsection we introduce the notion of an arrow and a 2-cell in \mathcal{C} representing an arrow and a 2-cell in $2\text{-}\mathcal{P}ro(\mathcal{C})$, and develop computational properties of 2-pro-objects which are necessary in our proof that the 2-category $2\text{-}\mathcal{P}ro(\mathcal{C})$ is closed under 2-cofiltered pseudo limits. In the third subsection we construct a 2-filtered category which serves as the index 2-category for the 2-filtered pseudolimit of 2-pro-objects (Definition 2.3.1 and Theorem 2.3.3). This is also inspired in a construction and proof for the same purpose found in [3], but which in our 2-dimensional case reveals itself very complex and difficult to manage effectively. We were forced to have recourse to this complicated construction because the conceptual treatment of this problem found in [1] does not apply in the 2-dimensional case. This is so because a 2-functor is not the pseudocolimit of 2-representables indexed by its 2-category of elements. Finally, in the last subsection we prove the universal properties of $2\text{-}\mathcal{P}ro(\mathcal{C})$ (Theorem 2.4.6), in a way which is novel even if applied to the classical theory of pro-objects.

1 Preliminaries on 2-categories

We distinguish between *small* and *large* sets. For us *legitimate* categories are categories with small hom sets, also called *locally small*. We freely consider without previous warning illegitimate categories with large hom sets, for example the category of all (legitimate) categories, or functor categories with large (legitimate) exponent. They are legitimate as categories in some higher universe, or they can be considered as convenient notational abbreviations for large collections of data. In fact, questions of size play no overt role in this paper, except that

we elect for simplicity to consider only small 2-pro-objects. We will explicitly mention whether the categories are legitimate or small when necessary. We reserve the notation Cat for the legitimate 2-category of small categories, and we will denote \mathcal{CAT} the illegitimate category (or 2-category) of all legitimate categories in some arbitrary sufficiently high universe.

Notation. 2-Categories will be denoted with the “mathcal” font \mathcal{C} , \mathcal{D} , ... , 2-functors with the capital “mathff” font, F , G , ... and 2-natural transformations, pseudonatural transformations and modifications with the greek alphabet. For objects in a 2-category, we will use capital “mathff” font C , D , ... , for arrows in a 2-category small case letters in “mathff” font f , g , ... , and for the 2-cells the greek alphabet. However, when a 2-category is intended to be used as the index 2-category of a 2-diagram, we will use small case letters i , j , ... to denote its objects, and small case letters u , v , ... to denote its arrows. Categories will be denoted with capital “mathff” font.

We begin with some background material on 2-categories. Most of this is standard, but some results (for which we provide proofs) do not appear to be in the literature. We also set notation and terminology as we will explicitly use in this paper.

1.1 Basic theory

Let Cat be the category of small categories. By a 2-category, we mean a Cat enriched category. A 2-functor, a 2-fully-faithful 2-functor, a 2-natural transformation and a 2-equivalence of 2-categories, are a Cat -functor, a Cat -fully-faithful functor, a Cat -natural transformation and a Cat -equivalence respectively.

In the sequel we will call *2-category* an structure satisfying the following descriptive definition free of the size restrictions implicit above. Given a 2-category, as usual, we denote horizontal composition by juxtaposition, and vertical composition by a “ \circ ”.

1.1.1. 2-Category. A 2-category \mathcal{C} consists on objects or 0-cells C , D ... , arrows or 1-cells f , g ... , and 2-cells α , β ,

$$\begin{array}{ccc} & f & \\ & \longrightarrow & \\ C & \xrightarrow{\alpha \Downarrow} & D \\ & g & \end{array}$$

The objects and the arrows form a category (called the underlying category of \mathcal{C}), with composition (called "horizontal") denoted by juxtaposition. For a fixed C and D , the arrows between them and the 2-cells between these arrows form a category $\mathcal{C}(C, D)$ under "vertical" composition, denoted by a " \circ ". There is also an associative horizontal composition between 2-cells denoted by juxtaposition, with units id_{id_C} . The following is the basic 2-category diagram:

$$\begin{array}{ccccc}
 & \xrightarrow{f} & & \xrightarrow{f'} & \\
 & \downarrow \alpha & & \downarrow \alpha' & \\
 C & \xrightarrow{g} & D & \xrightarrow{g'} & E \\
 & \downarrow \beta_h & & \downarrow \beta'_h & \\
 & \longrightarrow & & \longrightarrow &
 \end{array}$$

with the equations $(\beta' \beta) \circ (\alpha' \alpha) = (\beta' \circ \alpha')(\beta \circ \alpha)$, $id_{f'} id_f = id_{ff}$.

We consider juxtaposition more binding than " \circ ", thus $\alpha \beta \circ \gamma$ means $(\alpha \beta) \circ \gamma$. We will abuse notation by writing f instead of id_f for morphisms f and C instead of id_C for objects C .

1.1.2. Dual 2-Category. If \mathcal{C} is a 2-category, we denote by \mathcal{C}^{op} the 2-category with the same objects as \mathcal{C} but with $\mathcal{C}^{op}(C, D) = \mathcal{C}(D, C)$, i.e. we reverse the 1-cells but not the 2-cells.

1.1.3. 2-functor. A 2-functor $F : \mathcal{C} \rightarrow \mathcal{D}$ between 2-categories is an enriched functor over \mathcal{Cat} . As such, sends objects to objects, arrows to arrows and 2-cells to 2-cells, strictly preserving all the structure.

1.1.4. 2-fully-faithful. A 2-functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is said to be 2-fully-faithful if $\forall C, D \in \mathcal{C}$, $F_{C,D} : \mathcal{C}(C, D) \rightarrow \mathcal{D}(FC, FD)$ is an isomorphism of categories.

1.1.5. Pseudonatural. A pseudonatural transformation $\mathcal{C} \xrightarrow{\theta} \mathcal{D}$

between 2-functors consists in a family of morphisms $\{FC \xrightarrow{\theta_C} GC\}_{C \in \mathcal{C}}$ and a family of invertible 2-cells $\{Gf \theta_C \xrightarrow{\theta_f} \theta_D Ff\}_{C \xrightarrow{f} D \in \mathcal{C}}$

$$\begin{array}{ccc}
FC & \xrightarrow{\theta_C} & GC \\
Ff \downarrow & \Downarrow \theta_f & \downarrow Gf \\
FD & \xrightarrow{\theta_D} & GD
\end{array}$$

satisfying the following conditions:

$$PN0: \forall C \in \mathcal{C}, \quad \theta_{id_C} = id_{\theta_C}$$

$$PN1: \forall C \xrightarrow{f} D \xrightarrow{g} E, \quad \theta_{gf} = \theta_g Ff \circ Gg \theta_f.$$

$$PN2: \forall C \xrightarrow[\alpha \Downarrow]{f} D, \quad \theta_g \circ G\alpha \theta_C = \theta_D F\alpha \circ \theta_f$$

1.1.6. 2-Natural. A 2-natural transformation θ between 2-functors is a pseudonatural transformation such that $\theta_f = id \forall f \in \mathcal{C}$. Equivalently, it is a *Cat*-enriched natural transformation, that is, a natural transformation between the functors determined by F and G , such that for each

$$2\text{-cell } C \xrightarrow[\alpha \Downarrow]{f} D, \text{ the equation } G\alpha \theta_C = \theta_D F\alpha \text{ holds.} \quad \square$$

1.1.7. Modification. Given 2-functors F and G from \mathcal{C} to \mathcal{D} , a

$$\text{modification } F \xrightarrow[\rho \Downarrow]{\theta} G \text{ between pseudonatural transformations is a}$$

family $\{\theta_C \xrightarrow{\rho_C} \eta_C\}_{C \in \mathcal{C}}$ of 2-cells of \mathcal{D} such that:

$$\forall C \xrightarrow{f} D \in \mathcal{C}, \quad \rho_D Ff \circ \theta_f = \eta_f \circ Gf \rho_C.$$

As a particular case, we have modifications between 2-natural transformations, which are families of 2-cells as above satisfying $\rho_D Ff = Gf \rho_C$.

1.1.8. 2-Equivalence. A 2-functor $\mathcal{C} \xrightarrow{F} \mathcal{D}$ is said to be a 2-equivalence of 2-categories if there exists a 2-functor $\mathcal{D} \xrightarrow{G} \mathcal{C}$ and invertible 2-natural transformations $FG \xrightarrow{\alpha} id_{\mathcal{D}}$ and $GF \xrightarrow{\beta} id_{\mathcal{C}}$. G is said to be a quasi-inverse of F , and it is determined up to invertible 2-natural transformations.

1.1.9 Proposition. [11, 1.11] *A 2-functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is a 2-equivalence of 2-categories if and only if it is 2-fully-faithful and essentially surjective on objects.* \square

1.1.10. It is well known that 2-categories, 2-functors and 2-natural transformations form a 2-category (which actually underlies a 3-category) that we denote 2-CAT . Horizontal composition of 2-functors and vertical composition of 2-natural transformations are the usual ones, and the horizontal composition of 2-natural transformations is defined by:

$$\text{Given } \mathcal{C} \begin{array}{c} \xrightarrow{F} \\ \alpha \Downarrow \\ \xrightarrow{G} \end{array} \mathcal{D} \begin{array}{c} \xrightarrow{F'} \\ \alpha' \Downarrow \\ \xrightarrow{G'} \end{array} \mathcal{E}, \quad (\alpha' \alpha)_\mathcal{C} = \alpha'_{\mathcal{G}\mathcal{C}} \circ F'(\alpha_\mathcal{C}) \quad (= G'(\alpha_\mathcal{C}) \circ \alpha'_{\mathcal{F}\mathcal{C}}).$$

1.1.11 Definition. *Given two 2-categories \mathcal{C} and \mathcal{D} , we consider two 2-categories defined as follows:*

$\mathcal{H}om(\mathcal{C}, \mathcal{D})$: 2-functors and 2-natural transformations.

$\mathcal{H}om_p(\mathcal{C}, \mathcal{D})$: 2-functors and pseudonatural transformations.

In both cases the 2-cells are the modifications. To define compositions we draw the basic 2-category diagram:

$$\begin{array}{ccc} \xrightarrow{\theta} & \xrightarrow{\theta'} & \\ \Downarrow \rho & \Downarrow \rho' & \\ \mathcal{F} \xrightarrow{\eta} \mathcal{G} & \xrightarrow{\eta'} \mathcal{H} & \\ \Downarrow \varepsilon & \Downarrow \varepsilon' & \\ \xrightarrow{\mu} & \xrightarrow{\mu'} & \end{array} \quad \begin{array}{l} (\theta' \theta)_\mathcal{C} = \theta'_\mathcal{C} \theta_\mathcal{C} \\ (\rho' \rho)_\mathcal{C} = \rho'_\mathcal{C} \rho_\mathcal{C} \\ (\varepsilon \circ \rho)_\mathcal{C} = \varepsilon_\mathcal{C} \circ \rho_\mathcal{C} \end{array}$$

It is straightforward to check that these definitions determine a 2-category structure. \square

1.1.12 Remark. [9, I,4.2.] Evaluation determines a *quasifunctor* $\mathcal{H}om_p(\mathcal{C}, \mathcal{D}) \times \mathcal{C} \xrightarrow{ev} \mathcal{D}$ (in the sense of [9, I,4.1.], in particular, fixing a variable, it is a 2-functor in the other). In the strict case $\mathcal{H}om$, evaluation is actually a 2-bifunctor. \square

1.1.13 Remark. [9, I,4.2.] Both constructions $\mathcal{H}om$ and $\mathcal{H}om_p$ determine a bifunctor $2\text{-CAT}^{op} \times 2\text{-CAT} \rightarrow 2\text{-CAT}$. Given 2-functors

$\mathcal{C}' \xrightarrow{H_0} \mathcal{C}$ and $\mathcal{D} \xrightarrow{H_1} \mathcal{D}'$, and $F \xrightarrow[\eta]{\rho \Downarrow \theta} G$ in $\mathcal{H}om_\epsilon(\mathcal{C}, \mathcal{D})(F, G)$, the

definition $\mathcal{H}om_\epsilon(H_0, H_1)(F \xrightarrow[\eta]{\rho \Downarrow \theta} G) = H_1 F H_0 \xrightarrow[\begin{smallmatrix} H_1 \rho H_0 \Downarrow \\ H_1 \eta H_0 \end{smallmatrix}]{H_1 \theta H_0} H_1 G H_0$ determines a functor $\mathcal{H}om_\epsilon(\mathcal{C}, \mathcal{D})(F, G) \longrightarrow \mathcal{H}om_\epsilon(\mathcal{C}', \mathcal{D}')(H_1 F H_0, H_1 G H_0)$, and this assignation is bifunctorial in the variable $(\mathcal{C}, \mathcal{D})$ (here $\mathcal{H}om_\epsilon$ denotes either $\mathcal{H}om$ or $\mathcal{H}om_p$).

If \mathcal{C} and \mathcal{D} are 2-categories, the product 2-category $\mathcal{C} \times \mathcal{D}$ is constructed in the usual way, and this together with the 2-category $\mathcal{H}om(\mathcal{C}, \mathcal{D})$ determine a symmetric cartesian closed structure as follows (see [11, chapter 2] or [9, I,2.3.]):

1.1.14 Proposition. *The usual definitions determine an isomorphism of 2-categories :*

$$\mathcal{H}om(\mathcal{C}, \mathcal{H}om(\mathcal{D}, \mathcal{A})) \xrightarrow{\cong} \mathcal{H}om(\mathcal{C} \times \mathcal{D}, \mathcal{A}).$$

Composing with the symmetry $\mathcal{C} \times \mathcal{D} \xrightarrow{\cong} \mathcal{D} \times \mathcal{C}$ yields an isomorphism:

$$\mathcal{H}om(\mathcal{C}, \mathcal{H}om(\mathcal{D}, \mathcal{A})) \xrightarrow{\cong} \mathcal{H}om(\mathcal{D}, \mathcal{H}om(\mathcal{C}, \mathcal{A})). \quad \square$$

We use the following notation:

Notation: Let \mathcal{C} be a 2-category, $C \in \mathcal{C}$ and $D \xrightarrow[\mathbf{g}]{\mathbf{f}} E \in \mathcal{C}$.

1. $f_* : \mathcal{C}(C, D) \xrightarrow{f_*} \mathcal{C}(C, E)$, $f_*(h \xrightarrow{\beta} h') = (fh \xrightarrow{f\beta} fh')$.
2. $f^* : \mathcal{C}(E, C) \xrightarrow{f^*} \mathcal{C}(D, C)$, $f^*(h \xrightarrow{\beta} h') = (hf \xrightarrow{\beta f} hf')$.
3. $\alpha_* : f_* \xrightarrow{\alpha_*} g_*$, $(\alpha_*)_h = \alpha h$.
4. $\alpha^* : f^* \xrightarrow{\alpha^*} g^*$, $(\alpha^*)_h = h \alpha$.
5. $\mathcal{C} \xrightarrow{\mathcal{C}(C, -)} \mathcal{C}at : \mathcal{C}(C, -)(D \xrightarrow[\mathbf{g}]{\mathbf{f}} E) = (\mathcal{C}(C, D) \xrightarrow[\mathbf{g}^*]{\mathbf{f}_*} \mathcal{C}(C, E))$.

6. $\mathcal{C}^{op} \xrightarrow{\mathcal{C}(-, \mathcal{C})} \mathcal{C}at: \mathcal{C}(-, \mathcal{C})(D \xrightarrow[\underset{\mathbf{g}}{\alpha \Downarrow}}{\mathbf{f}}]{\mathbf{f}} E) = (\mathcal{C}(D, \mathcal{C}) \xrightarrow[\underset{\mathbf{g}^*}{\alpha^* \Downarrow}]{\mathbf{f}^*} \mathcal{C}(E, \mathcal{C}))$.
7. We will also denote by \mathbf{f}^* the 2-natural transformation from $\mathcal{C}(E, -)$ to $\mathcal{C}(D, -)$ defined by $(\mathbf{f}^*)_C = \mathbf{f}^*$.
8. We will also denote by \mathbf{f}_* the 2-natural transformation from $\mathcal{C}(-, D)$ to $\mathcal{C}(-, E)$ defined by $(\mathbf{f}_*)_C = \mathbf{f}_*$.
9. We will also denote by α^* the modification from \mathbf{f}^* to \mathbf{g}^* defined by $(\alpha^*)_C = \alpha^*$.
10. We will also denote by α_* the modification from \mathbf{f}_* to \mathbf{g}_* defined by $(\alpha_*)_C = \alpha_*$. \square

1.1.15. *Given a locally small 2-category \mathcal{C} , the Yoneda 2-functors are the following (note that each one is the other for the dual 2-category):*

- a. $\mathcal{C} \xrightarrow{y^{(-)}} \mathcal{H}om(\mathcal{C}, \mathcal{C}at)^{op}$, $y^C = \mathcal{C}(C, -)$, $y^f = \mathbf{f}^*$, $y^\alpha = \alpha^*$.
- b. $\mathcal{C} \xrightarrow{y^{(-)}} \mathcal{H}om(\mathcal{C}^{op}, \mathcal{C}at)$, $y_C = \mathcal{C}(-, C)$, $y_f = \mathbf{f}_*$, $y_\alpha = \alpha_*$.

Recall the Yoneda Lemma for enriched categories over $\mathcal{C}at$. We consider explicitly only the case a. in 1.1.15.

1.1.16 Proposition (Yoneda lemma). *Given a locally small 2-category \mathcal{C} , a 2-functor $F: \mathcal{C} \rightarrow \mathcal{C}at$ and an object $C \in \mathcal{C}$, there is an isomorphism of categories, natural in F .*

$$\begin{array}{ccc} \mathcal{H}om(\mathcal{C}, \mathcal{C}at)(\mathcal{C}(C, -), F) & \xrightarrow{h} & FC \\ \theta \xrightarrow{\rho} \eta & \longmapsto & \theta_C(id_C) \xrightarrow{(\rho_C)id_C} \eta_C(id_C) \end{array}$$

Proof. The application h has an inverse

$$\begin{array}{ccc} FC & \xrightarrow{\ell} & \mathcal{H}om(\mathcal{C}, \mathcal{C}at)(\mathcal{C}(C, -), F) \\ C \xrightarrow{f} D & \longmapsto & \ell C \xrightarrow{\ell f} \ell D \end{array}$$

where $(\ell C)_D(f \xrightarrow{\alpha} g) = \mathbf{F}f(C) \xrightarrow{(\mathbf{F}\alpha)_C} \mathbf{F}g(C)$ and $((\ell f)_D)_f = \mathbf{F}f$. \square

1.1.17 Corollary. *The Yoneda 2-functors in 1.1.15 are 2-fully-faithful.* \square

Beyond the theory of Cat -enriched categories, the lemma also holds for pseudonatural transformations in the following way:

1.1.18 Proposition (Pseudo Yoneda lemma). *Given a locally small 2-category \mathcal{C} , a 2-functor $F : \mathcal{C} \rightarrow Cat$ and an object $C \in \mathcal{C}$, there is an equivalence of categories, natural in F .*

$$\begin{array}{ccc} \mathcal{H}om_p(\mathcal{C}, Cat)(\mathcal{C}(C, -), F) & \xrightarrow{\tilde{h}} & FC \\ \theta \xrightarrow{\rho} \eta & \longmapsto & \theta_C(id_C) \xrightarrow{(\rho_C)id_C} \eta_C(id_C) \end{array}$$

Furthermore, the quasi-inverse $\tilde{\ell}$ is a section of \tilde{h} , $\tilde{h}\tilde{\ell} = id$.

Proof. \tilde{h} and $\tilde{\ell}$ are defined as in 1.1.16, but now $\tilde{\ell}$ is only a section quasi-inverse of \tilde{h} . The details can be checked by the reader. One can find a guide in [13] for the case of lax functors and bicategories. We refer to the arguing and the notation there: In our case, the unit η is the equality because F is a 2-functor, and the counit ϵ is an isomorphism because α is pseudonatural and the unitor r is the equality. \square

1.1.19 Corollary. *For any locally small 2-category \mathcal{C} , and $C \in \mathcal{C}$, the inclusion $\mathcal{H}om(\mathcal{C}, Cat)(\mathcal{C}(C, -), F) \xrightarrow{i} \mathcal{H}om_p(\mathcal{C}, Cat)(\mathcal{C}(C, -), F)$ has a retraction α , natural in F , $\alpha i = id$, $i\alpha \cong id$, which determines an equivalence of categories.*

Proof. Note that $i = \tilde{\ell}h$, then define $\alpha = \ell\tilde{h}$. \square

1.1.20 Corollary. *The Yoneda 2-functors in 1.1.15 can be considered as 2-functors landing in the $\mathcal{H}om_p$ 2-functor categories. In this case, they are pseudo-fully faithful (meaning that they determine equivalences and not isomorphisms between the hom categories).* \square

1.2 Weak limits and colimits

By *weak* we understand any of the several ways universal properties can be relaxed in 2-categories. Note that pseudolimits and pseudocolimits (already considered in [2]) require isomorphisms, and have many advantages over bilimits and bicolimits, which only require equivalences. Their universal properties are both stronger and more convenient to use, and they play the principal role in this paper. The defining universal properties characterize bilimits up to equivalence and pseudolimits up to isomorphism.

Notation We consider pseudocolimits $\varinjlim_{i \in \mathcal{I}} Fi$, and bicolimits $\mathbf{bi}\varinjlim_{i \in \mathcal{I}} Fi$, of covariant 2-functors, and its dual concepts, pseudolimits $\varprojlim_{i \in \mathcal{I}} Fi$, and bilimits $\mathbf{bi}\varprojlim_{i \in \mathcal{I}} Fi$, of contravariant 2-functors.

1.2.1 Definition. Let $F : \mathcal{I} \rightarrow \mathcal{A}$ be a 2-functor and A an object of \mathcal{A} . A pseudocone for F with vertex A is a pseudonatural transformation from F to the 2-functor which is constant at A , i.e. it consists in a family of morphisms of \mathcal{A} $\{Fi \xrightarrow{\theta_i} A\}_{i \in \mathcal{I}}$ and a family of invertible 2-cells of \mathcal{A} $\{\theta_i \xrightarrow{\theta_u} \theta_j Fu\}_{i \xrightarrow{u} j \in \mathcal{I}}$ satisfying the following equations:

$$\begin{aligned} PC0: & \theta_{id_i} = id_{\theta_i} . \\ PC1: & \forall i \xrightarrow{u} j \xrightarrow{v} k \in \mathcal{I}, \quad \theta_v Fu \circ \theta_u = \theta_{vu} \\ PC2: & \forall i \xrightarrow[\alpha \Downarrow]{u} j \in \mathcal{I}, \quad \theta_i = \theta_j F\alpha \circ \theta_u \end{aligned}$$

A morphism of pseudocones between θ and η with the same vertex is a modification, i.e. a family of 2-cells of \mathcal{A} $\{\theta_i \xrightarrow{\rho_i} \eta_i\}_{i \in \mathcal{I}}$ satisfying the following equation:

$$PCM: \quad \eta_u \circ \rho_i = \rho_j Fu \circ \theta_u .$$

Pseudocones form a category $PC_{\mathcal{A}}(F, A) = \mathcal{H}om_p(\mathcal{I}, \mathcal{A})(F, A)$ furnished with a pseudocone $PC_{\mathcal{A}}(F, A) \rightarrow \mathcal{A}(Fi, A)$, $\{\theta_i\}_{i \in \mathcal{I}} \mapsto \theta_i$, for the 2-functor $\mathcal{I}^{op} \xrightarrow{\mathcal{A}(F(-), A)} \mathcal{CAT}$.

1.2.2 Remark.

Since $\mathcal{H}om_p(\mathcal{I}, \mathcal{A})$ is a 2-category, it follows:

a. Pseudocones determine a 2-bifunctor $\mathcal{H}om(\mathcal{I}, \mathcal{A})^{op} \times \mathcal{A} \xrightarrow{\text{PC}_{\mathcal{A}}} \mathcal{CAT}$.

From Remark 1.1.13 it follows in particular:

b. A 2-functor $\mathcal{A} \xrightarrow{\text{H}} \mathcal{B}$ induces a functor between the categories of pseudocones $\text{PC}_{\mathcal{A}}(\text{F}, \text{A}) \xrightarrow{\text{PC}_{\text{H}}} \text{PC}_{\mathcal{B}}(\text{HF}, \text{HA})$. \square

1.2.3 Definition. *The pseudocolimit in \mathcal{A} of the 2-functor F is the universal pseudocone, denoted $\{\text{Fi} \xrightarrow{\lambda_i} \varinjlim_{i \in \mathcal{I}} \text{Fi}\}_{i \in \mathcal{I}}$, in the sense that $\forall \text{A} \in \mathcal{A}$, pre-composition with the λ_i is an isomorphism of categories $\mathcal{A}(\varinjlim_{i \in \mathcal{I}} \text{Fi}, \text{A}) \xrightarrow{\lambda^*} \text{PC}_{\mathcal{A}}(\text{F}, \text{A})$. Equivalently, there is an isomorphism of categories $\mathcal{A}(\varinjlim_{i \in \mathcal{I}} \text{Fi}, \text{A}) \xrightarrow{\cong} \varprojlim_{i \in \mathcal{I}^{op}} \mathcal{A}(\text{Fi}, \text{A})$ commuting with the pseudocones. Remark that there is also an isomorphism of categories $\text{PC}_{\mathcal{A}}(\text{F}, \text{A}) \xrightarrow{\cong} \varprojlim_{i \in \mathcal{I}^{op}} \mathcal{A}(\text{Fi}, \text{A})$*

Requiring λ^ to be an equivalence (which implies that also the other two isomorphisms above are equivalences) defines the notion of bicolimit. Clearly, pseudocolimits are bicolimits.*

We omit the explicit consideration of the dual concepts. \square

It is well known that in the strict 2-functor 2-categories the strict limits and colimits are performed pointwise (if they exists in the codomain category). Here we establish this fact for the pseudo limits and pseudocolimits in both the strict and the pseudo 2-functor 2-categories. Abusing notation we can say that the formula $(\varinjlim_{i \in \mathcal{I}} \text{Fi})(\text{C}) = \varinjlim_{i \in \mathcal{I}} \text{Fi}(\text{C})$ holds in both 2-categories. The verification of this is straightforward but requires some care.

1.2.4 Proposition. *Let $\mathcal{I} \xrightarrow{\text{F}} \mathcal{A}$, $i \mapsto \text{Fi}$ be a 2-functor where \mathcal{A} is either $\mathcal{H}om(\mathcal{C}, \mathcal{D})$ or $\mathcal{H}om_p(\mathcal{C}, \mathcal{D})$. For each $\text{C} \in \mathcal{C}$ let $\text{FiC} \xrightarrow{\lambda_i^{\text{C}}} \text{LC}$ be a pseudocolimit pseudocone in \mathcal{D} for the 2 functor $\mathcal{I} \xrightarrow{\text{F}} \mathcal{A} \xrightarrow{\text{ev}(-, \text{C})} \mathcal{D}$ (where ev is evaluation, see 1.1.12). Then LC is 2-functorial in \mathcal{C} in such a way that λ_i^{C} becomes 2-natural and $\text{Fi} \xrightarrow{\lambda_i} \text{L}$ is a pseudocolimit pseudocone in \mathcal{A} in both cases. By duality the same assertion holds for pseudolimits.*

Proof. Given $C \xrightarrow[\alpha \Downarrow]{f} D$ in \mathcal{C} , evaluation determines a 2-cell in $\mathcal{H}om(\mathcal{I}, \mathcal{D})$ $FC \xrightarrow[\text{Fg}]{\text{Ff}} FD = ev(F(-), C \xrightarrow[\alpha \Downarrow]{f} D)$. (note that $(FC)_i = F_i C$, and similarly for f , g and α). Then, for each $X \in \mathcal{D}$, it follows (from Remark 1.2.2 a.) that precomposing with this 2-cell determines a 2-cell (clearly 2-natural in the variable X) in the right leg of the diagram below. Since the rows are isomorphisms, there is a unique 2-cell (also natural in the variable X) in the left leg which makes the diagram commutative.

$$\begin{array}{ccc} \mathcal{D}(\text{LD}, X) & \xrightarrow[\cong]{(\lambda^D)^*} & \text{PC}_{\mathcal{D}}(\text{FD}, X) \\ \downarrow \Rightarrow \downarrow & & \downarrow \Rightarrow \downarrow \\ \mathcal{D}(\text{LC}, X) & \xrightarrow[\cong]{(\lambda^C)^*} & \text{PC}_{\mathcal{D}}(\text{FC}, X) \end{array}$$

Then, by the Yoneda lemma 1.1.17, the left leg is given by precomposing with a unique 2-cell in \mathcal{D} , that we denote $\text{LC} \xrightarrow[\text{Lg}]{\text{Lf}} \text{LD}$. It is clear by uniqueness that this determines a 2-functor $\mathcal{C} \xrightarrow{\text{L}} \mathcal{D}$.

Putting $X = \text{LD}$ in the upper left corner and tracing the identity down the diagram yields the following commutative diagram of pseudocones in \mathcal{D} :

$$\begin{array}{ccc} F_i C & \xrightarrow{\lambda_i^C} & \text{LC} \\ F_i f \downarrow \begin{array}{c} F_i \alpha \\ \Rightarrow \\ F_i g \end{array} & & \downarrow \begin{array}{c} \text{L} \alpha \\ \Rightarrow \\ \text{L} g \end{array} \\ F_i D & \xrightarrow{\lambda_i^D} & \text{LD} \end{array}$$

This shows that L is furnished with a pseudocone for F and that the λ_i are 2-natural. It only remains to check the universal property:

Let $\mathcal{C} \xrightarrow{\text{G}} \mathcal{D}$ be a 2-functor, consider the 2-functor $\mathcal{A} \xrightarrow{ev(-, \mathcal{C})} \mathcal{D}$. We have the following diagram, where the right leg is given by

Remark 1.2.2 b.:

$$\begin{array}{ccc} \mathcal{A}(\mathbf{L}, \mathbf{G}) & \xrightarrow{\lambda^*} & \mathrm{PC}_{\mathcal{A}}(\mathbf{F}, \mathbf{G}) \\ \downarrow \mathrm{ev}(-, \mathbf{C}) & & \downarrow \mathrm{PC}_{\mathrm{ev}(-, \mathbf{C})} \\ \mathcal{D}(\mathrm{LC}, \mathrm{GC}) & \xrightarrow[\cong]{(\lambda^{\mathbf{C}})^*} & \mathrm{PC}_{\mathcal{D}}(\mathrm{FC}, \mathrm{GC}) \end{array}$$

We prove now that the upper row is an isomorphism. Given $\mathbf{F}_i \xrightarrow[\eta_i]{\rho_i \downarrow} \mathbf{G}$ in $\mathrm{PC}_{\mathcal{A}}(\mathbf{F}, \mathbf{G})$, it follows there exists a unique $\mathrm{LC} \xrightarrow[\tilde{\eta}\mathbf{C}]{\tilde{\rho}\mathbf{C} \downarrow} \mathrm{GC}$ in $\mathcal{D}(\mathrm{LC}, \mathrm{GC})$ such that $\tilde{\rho}\mathbf{C} \lambda_i^{\mathbf{C}} = \rho_i \mathbf{C}$. It is necessary to show that this 2-cell actually lives in \mathcal{A} . This has to be checked for any $\mathbf{C} \xrightarrow[\mathbf{g}]{\alpha \downarrow} \mathbf{D}$ in \mathcal{C} . In both cases it can be done considering the isomorphism $\mathcal{D}(\mathrm{LC}, \mathrm{GD}) \xrightarrow[\cong]{(\lambda^{\mathbf{C}})^*} \mathrm{PC}_{\mathcal{D}}(\mathrm{FC}, \mathrm{GD})$. \square

We precise now what we do consider as *preservation* properties of a 2-functor. We do it in the case of pseudolimits and bilimits, but the same clearly applies to pseudocolimits and bicolimits. Let $\mathcal{I}^{op} \xrightarrow{\mathbf{X}} \mathcal{C} \xrightarrow{\mathbf{H}} \mathcal{A}$ be any 2-functors.

1.2.5 Definition. *We say that \mathbf{H} preserves a pseudolimit (resp. bilimit) pseudocone $\mathbf{L} \xrightarrow{\pi_i} \mathbf{X}_i$ in \mathcal{C} , if $\mathbf{H}\mathbf{L} \xrightarrow{\mathbf{H}\pi_i} \mathbf{H}\mathbf{X}_i$ is a pseudolimit (resp. bilimit) pseudocone in \mathcal{A} . Equivalently, if the (usual) comparison arrow is an isomorphism (resp. an equivalence) in \mathcal{A} .*

Note that by the very definition, the 2-representable 2-functors preserve pseudolimits and bilimits. Also, from proposition 1.2.4 it follows:

1.2.6 Proposition. *The Yoneda 2-functors in 1.1.15 preserve pseudolimits.* \square

Recall that small pseudolimits and pseudocolimits of locally small categories exist and are locally small, as well that the 2-category $\mathcal{C}at$ of small categories has all small pseudolimits and pseudocolimits (see for example [4], [12]).

1.2.7. We refer to the explicit construction of pseudolimits of category valued 2-functors, which is similar to the construction of pseudolimits of category-valued functors in [2, Exposé VI 6.], see full details in [5].

It is also key to our work the explicit construction of 2-filtered pseudocolimits of category valued 2-functors developed in [8]. We recall this now.

1.2.8 Definition (Kennison, [10]). *Let \mathcal{C} be a 2-category. \mathcal{C} is said to be 2-filtered if the following axioms are satisfied:*

F0. Given two objects $C, D \in \mathcal{C}$, there exists an object $E \in \mathcal{C}$ and arrows $C \rightarrow E, D \rightarrow E$.

F1. Given two arrows $C \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} D$, there exists an arrow $D \xrightarrow{h} E$ and an invertible 2-cell $\alpha : hf \cong hg$.

F2. Given two 2-cells $C \begin{array}{c} \xrightarrow{\alpha \Downarrow \beta \Downarrow} \\ \xrightarrow{\alpha \Downarrow \beta \Downarrow} \end{array} D$ there exists an arrow $D \xrightarrow{h} E$ such that $h\alpha = h\beta$.

The dual notion of 2-cofiltered 2-category is given by the duals of axioms F0, F1 and F2.

1.2.9. Construction LL (Dubuc-Street [8]) Let \mathcal{I} be a 2-filtered 2-category and $F : \mathcal{I} \rightarrow \mathcal{Cat}$ a 2-functor. We define a category $\mathcal{L}(F)$ in two steps as follows:

First step ([8, Definition 1.5]):

Objects: (C, i) with $C \in Fi$

Premorphisms: A premorphism between (C, i) and (D, j) is a triple (u, f, v) where $i \xrightarrow{u} k, j \xrightarrow{v} k$ in \mathcal{I} and $F(u)(C) \xrightarrow{f} F(v)(D)$ in Fk .

Homotopies: An homotopy between two premorphisms (u_1, f_1, v_1) and (u_2, f_2, v_2) is a quadruple $(w_1, w_2, \alpha, \beta)$ where $k_1 \xrightarrow{w_1} k, k_2 \xrightarrow{w_2} k$ are 1-cells of \mathcal{I} and $w_1v_1 \xrightarrow{\alpha} w_2v_2, w_1u_1 \xrightarrow{\beta} w_2u_2$ are invertible 2-cells of \mathcal{I} such that the following diagram commutes in Fk :

$$\begin{array}{ccc} F(w_1)F(u_1)(C) = F(w_1u_1)(C) & \xrightarrow{F(\beta)C} & F(w_2u_2)(C) = F(w_2)F(u_2)(C) \\ \downarrow F(w_1)(f_1) & & \downarrow F(w_2)(f_2) \\ F(w_1)F(v_1)(D) = F(w_1v_1)(D) & \xrightarrow{F(\alpha)D} & F(w_2v_2)(D) = F(w_2)F(v_2)(D) \end{array}$$

We say that two premorphisms f_1, f_2 are equivalent if there is an homotopy between them. In that case, we write $f_1 \sim f_2$.

Equivalence is indeed an equivalence relation, and premorphisms can be (non uniquely) composed. Up to equivalence, composition is independent of the choice of representatives and of the choice of the composition between them. Since associativity holds and identities exist, the following actually does define a category:

Second step ([8, Definition 1.13]):

Objects: (C, i) with $C \in Fi$.

Morphisms: equivalence classes of premorphisms.

Composition: defined by composing representative premorphisms.

1.2.10 Proposition. [8, Theorem 1.19] *Let \mathcal{I} be a 2-filtered 2-category, $F : \mathcal{I} \rightarrow \mathcal{C}at$ a 2-functor, $i \xrightarrow{u} j$ in \mathcal{I} and $C \xrightarrow{f} D \in Fi$. The following formulas define a pseudocone $F \xrightarrow{\lambda} \mathcal{L}(F)$:*

$$\lambda_i(C) = (C, i) \quad \lambda_i(f) = [i, f, i] \quad (\lambda_u)_C = [u, Fu(C), j]$$

which is a pseudocolimit for the 2-functor F . □

1.3 Further results.

A. Joyal pointed to us the notion of *flexible* functors, related with some of our results on pseudo colimits of representable 2-functors. We recall now this notion since it bears some significance for the concept of 2-pro-object developed in this paper. Any 2-pro-object determines a 2-functor which is flexible, and some of our results find their right place stated in the context of flexible 2-functors.

Warning: *In this subsection 2-categories are assumed to be locally small, except the illegitimate constructions $\mathcal{H}om$ and $\mathcal{H}om_p$.*

The inclusion $\mathcal{H}om(\mathcal{C}, \mathcal{C}at) \xrightarrow{i} \mathcal{H}om_p(\mathcal{C}, \mathcal{C}at)$ has a left adjoint $(-)' \dashv i$, we refer the reader to [4]. The 2-natural counit of this adjunction $F' \xrightarrow{\varepsilon_F} F$ is an equivalence in $\mathcal{H}om_p(\mathcal{C}, \mathcal{C}at)$, with a section given by the pseudonatural unit $F \xrightarrow{\eta_F} F'$, $\varepsilon_F \eta_F = id_F$, $\eta_F \varepsilon_F \cong id_{F'}$, [4, Proposition 4.1.]

1.3.1 Definition. [4, Proposition 4.2] A 2-functor $\mathcal{C} \xrightarrow{F} \mathcal{C}at$ is flexible if the counit $F' \xrightarrow{\varepsilon_F} F$ has a 2-natural section $F \xrightarrow{\lambda} F'$, $\varepsilon_F \lambda = id_F$, $\lambda \varepsilon_F \cong id_{F'}$, which determines an equivalence in $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)$.

We state now a useful characterization of flexible 2-functors F independent of the left adjoint $(-)'$, the proof will appear elsewhere [6].

1.3.2 Proposition. A 2-functor $\mathcal{C} \xrightarrow{F} \mathcal{C}at$ is flexible \iff for all 2-functors G , the inclusion $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)(F, G) \xrightarrow{i_G} \mathcal{H}om_p(\mathcal{C}, \mathcal{C}at)(F, G)$ has a retraction α_G natural in G , $\alpha_G i_G = id$, $i_G \alpha_G \cong id$, which determines an equivalence of categories. \square

Let $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)_f$ and $\mathcal{H}om_p(\mathcal{C}, \mathcal{C}at)_f$ be the subcategories whose objects are the flexible 2-functors. We have the following corollaries:

1.3.3 Corollary. The 2-categories $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)_f$ and $\mathcal{H}om_p(\mathcal{C}, \mathcal{C}at)_f$ are pseudoequivalent in the sense they have the same objects and retract equivalent hom categories. \square

We mention that following the usual lines (based in the axiom of choice) in the proof of 1.1.9, it can be seen that the inclusion 2-functor $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)_f \rightarrow \mathcal{H}om_p(\mathcal{C}, \mathcal{C}at)_f$ has the identity (on objects) as a retraction quasi-inverse *pseudofunctor*, with the equality as the invertible pseudonatural transformation $F \xrightarrow{=} F$ in $\mathcal{H}om_p(\mathcal{C}, \mathcal{C}at)_f$.

An important property of flexible 2-functors, false in general, is the following:

1.3.4 Corollary. Let $\theta : G \Rightarrow F \in \mathcal{H}om(\mathcal{C}, \mathcal{C}at)_f$ be such that $\theta_C : GC \rightarrow FC$ is an equivalence of categories for each $C \in \mathcal{C}$. Then, θ is an equivalence in $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)_f$.

Proof. It is easy to check that there is a pseudonatural transformation $\eta' : F \Rightarrow G$ such that $\theta \eta' \cong F$ and $\eta' \theta \cong G$ in $\mathcal{H}om_p(F, F)$ and $\mathcal{H}om_p(G, G)$ respectively. Now, by 1.3.2, there is a 2-natural transformation $\eta : F \Rightarrow G$ such that $\eta \cong \eta'$ in $\mathcal{H}om_p(F, G)$. Then, $\theta \eta \cong F$ and $\eta \theta \cong G$ in $\mathcal{H}om(F, F)$ and $\mathcal{H}om(G, G)$ respectively and so θ is an equivalence in $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)$. \square

1.3.5 Proposition. Small pseudocolimits of flexible 2-functors are flexible.

Proof. Let $F = \varinjlim_{j \in \mathcal{I}} F_j$, where each F_j is flexible, and let G be any other 2-functor. Set $\mathcal{A} = \mathcal{H}om(\mathcal{C}, \mathcal{C}at)$ and $\mathcal{A}_p = \mathcal{H}om_p(\mathcal{C}, \mathcal{C}at)$. Then:

$$\mathcal{A}(F, G) \cong \varprojlim_{j \in \mathcal{I}} \mathcal{A}(F_j, G) \xrightarrow{i} \varprojlim_{j \in \mathcal{I}} \mathcal{A}_p(F_j, G) \cong \mathcal{A}_p(F, G).$$

The two isomorphisms are given by definition 1.2.3. The arrow i is the pseudolimit of the equivalences with retraction quasi-inverses corresponding to each F_j . It is not difficult to check that i is also such an equivalence. \square

It follows also from 1.3.2 that the pseudo-Yoneda lemma (1.1.18, 1.1.19) says that the representable 2-functors are flexible, so we have:

1.3.6 Corollary. *Small pseudocolimits of representable 2-functors are flexible.* \square

Note that 1.3.5 and 1.3.6 hold for any pseudocolimit that may exist.

2 2-Pro-objects

Warning: *In this section 2-categories are assumed to be locally small, except illegitimate constructions as $\mathcal{H}om$, $\mathcal{H}om_p$ or 2-CAT .*

The main results of this paper are in this section. In the first subsection we define the 2-category of 2-pro-objects of a 2-category \mathcal{C} and establish the basic formula for the morphisms and 2-cells of this 2-category. Then in the next subsection we develop the notion of a 2-cell in \mathcal{C} representing a 2-cell in $2\text{-Pro}(\mathcal{C})$, inspired in the 1-dimensional notion of an arrow representing a morphism of pro-objects found in [3]. We use this in the third subsection to construct the 2-filtered category which serves as the index 2-category for the 2-filtered pseudolimit of 2-pro-objects. This is also inspired in a construction for the same purpose found in [3]. We were forced to have recourse to this complicated construction because the conceptual treatment of this problem found in [1] does not apply in the 2-category case. This is so because a 2-functor is not the pseudocolimit of 2-representables indexed by its 2-category of elements. Finally, in the last subsection we prove the universal properties of $2\text{-Pro}(\mathcal{C})$.

2.1 Definition of the 2-category of 2-pro-objects

In this subsection we define the 2-category of 2-pro-objects of a fixed 2-category and prove its basic properties. A 2-pro-object over a 2-category \mathcal{C} will be a small 2-cofiltered diagram in \mathcal{C} and it will be the pseudolimit of it's own diagram in the 2-category $2\text{-Pro}(\mathcal{C})$.

2.1.1 Definition. *Let \mathcal{C} be a 2-category. We define the 2-category of 2-pro-objects of \mathcal{C} , which we denote by $2\text{-Pro}(\mathcal{C})$, as follows:*

1. *Its objects are the 2-functors $\mathcal{I}^{op} \xrightarrow{\mathbf{X}} \mathcal{C}$, $\mathbf{X} = (\mathbf{X}_i, \mathbf{X}_u, \mathbf{X}_\alpha)_{i, u, \alpha \in \mathcal{I}}$, with \mathcal{I} a small 2-filtered 2-category. Often we are going to abuse the notation by saying $\mathbf{X} = (\mathbf{X}_i)_{i \in \mathcal{I}}$.*
2. *If $\mathbf{X} = (\mathbf{X}_i)_{i \in \mathcal{I}}$ and $\mathbf{Y} = (\mathbf{Y}_j)_{j \in \mathcal{J}}$ are two 2-pro-objects,*

$$\begin{aligned} 2\text{-Pro}(\mathcal{C})(\mathbf{X}, \mathbf{Y}) &= \mathcal{H}om(\mathcal{C}, \mathcal{C}at)^{op}(\varprojlim_{i \in \mathcal{I}^{op}} \mathcal{C}(\mathbf{X}_i, -), \varprojlim_{j \in \mathcal{J}^{op}} \mathcal{C}(\mathbf{Y}_j, -)) \\ &= \mathcal{H}om(\mathcal{C}, \mathcal{C}at)(\varinjlim_{j \in \mathcal{J}} \mathcal{C}(\mathbf{Y}_j, -), \varinjlim_{i \in \mathcal{I}} \mathcal{C}(\mathbf{X}_i, -)) \end{aligned}$$

The compositions are given by the corresponding compositions in the 2-category $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)^{op}$ so it is easy to check that $2\text{-Pro}(\mathcal{C})$ is indeed a 2-category.

2.1.2 Proposition. *By definition there is a 2-fully-faithful 2-functor $2\text{-Pro}(\mathcal{C}) \xrightarrow{\mathbf{L}} \mathcal{H}om(\mathcal{C}, \mathcal{C}at)^{op}$. Thus, there is a contravariant 2-equivalence of 2-categories $2\text{-Pro}(\mathcal{C}) \xrightarrow{\mathbf{L}} \mathcal{H}om(\mathcal{C}, \mathcal{C}at)_{fc}^{op}$, where $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)_{fc}$ stands for the full subcategory of $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)$ whose objects are those 2-functors which are small 2-filtered pseudocolimits of representable 2-functors. However, it is important to note that this equivalence is not injective on objects. \square*

From Corollary 1.3.6 it follows:

2.1.3 Proposition. *For any 2-pro-object \mathbf{X} , the corresponding 2-functor $\mathbf{L}\mathbf{X}$ is flexible. \square*

2.1.4 Remark. If we use pseudonatural transformations to define morphisms of 2-pro-objects we obtain a 2-category $2\text{-}\mathcal{P}ro_p(\mathcal{C})$, which anyway, by 2.1.3, results pseudoequivalent (see 1.3.3) to $2\text{-}\mathcal{P}ro(\mathcal{C})$, with the same objects and retract equivalent hom categories. We think our choice of morphisms, which is much more convenient to use, will prove to be the good one for the applications.

Next we establish the basic formula which is essential in many computations in the 2-category $2\text{-}\mathcal{P}ro(\mathcal{C})$:

2.1.5 Proposition. *There is an isomorphism of categories:*

$$(2.1.5) \quad 2\text{-}\mathcal{P}ro(\mathcal{C})(\mathbf{X}, \mathbf{Y}) \cong \underset{j \in \mathcal{J}^{op}}{\leftarrow \text{Lim}} \underset{i \in \mathcal{I}}{\text{Lim}} \mathcal{C}(X_i, Y_j)$$

Proof.

$$\begin{aligned} 2\text{-}\mathcal{P}ro(\mathcal{C})(\mathbf{X}, \mathbf{Y}) &= \text{Hom}(\mathcal{C}, \text{Cat})(\underset{j \in \mathcal{J}}{\text{Lim}} \mathcal{C}(Y_j, -), \underset{i \in \mathcal{I}}{\text{Lim}} \mathcal{C}(X_i, -)) \cong \\ &\underset{j \in \mathcal{J}^{op}}{\leftarrow \text{Lim}} \text{Hom}(\mathcal{C}, \text{Cat})(\mathcal{C}(Y_j, -), \underset{i \in \mathcal{I}}{\text{Lim}} \mathcal{C}(X_i, -)) \cong \underset{j \in \mathcal{J}^{op}}{\leftarrow \text{Lim}} \underset{i \in \mathcal{I}}{\text{Lim}} \mathcal{C}(X_i, Y_j) \end{aligned}$$

The first isomorphism is due to 1.2.3 and the second one to 1.1.16. \square

2.1.6 Corollary. *The 2-category $2\text{-}\mathcal{P}ro(\mathcal{C})$ is locally small.*

2.1.7 Corollary. *There is a canonical 2-fully-faithful 2-functor $\mathcal{C} \xrightarrow{c} 2\text{-}\mathcal{P}ro(\mathcal{C})$ which sends an object of \mathcal{C} into the corresponding 2-pro-object with index 2-category $\{*\}$. Since this 2-functor is also injective on objects, we can identify \mathcal{C} with a 2-full subcategory of $2\text{-}\mathcal{P}ro(\mathcal{C})$. \square*

Where there is no risk of confusion, we will omit to indicate notationally this identification. By the very definition of $2\text{-}\mathcal{P}ro(\mathcal{C})$ it follows:

2.1.8 Proposition. *If $\mathbf{X} = (X_i)_{i \in \mathcal{I}}$ is any 2-pro-object of \mathcal{C} , then $\mathbf{X} = \underset{i \in \mathcal{I}^{op}}{\leftarrow \text{Lim}} X_i$ in $2\text{-}\mathcal{P}ro(\mathcal{C})$. \mathbf{X} is equipped with projections, for each $i \in \mathcal{I}$, $\mathbf{X} \xrightarrow{\pi_i} X_i$, and a pseudocone structure, for each $i \xrightarrow{u} j \in \mathcal{I}$, invertible 2-cells $\pi_i \xrightarrow{\pi_u} X_u \pi_j$.*

Under the isomorphism $2\text{-Pro}(\mathcal{C})(X, X_i) \cong \varinjlim_{k \in \mathcal{I}} \mathcal{C}(X_k, X_i)$ (2.1.5),

the projections $X \xrightarrow{\pi_i} X_i$ correspond to the object (id_{X_i}, i) in construction 1.2.9. \square

Note that from this proposition it follows:

2.1.9 Remark. Given any two pro-objects $X, Z \in 2\text{-Pro}(\mathcal{C})$, there is an isomorphism of categories $2\text{-Pro}(\mathcal{C})(Z, X) \xrightarrow{\cong} \text{PC}_{2\text{-Pro}(\mathcal{C})}(Z, cX)$, where $\text{PC}_{2\text{-Pro}(\mathcal{C})}$ is the category of pseudocones for the 2-functor cX with vertex Z .

It is important to note that when $\varprojlim_{i \in \mathcal{I}^{op}} X_i$ exists in \mathcal{C} , this pseudolimit would not be isomorphic to X in $2\text{-Pro}(\mathcal{C})$. In general, the functor c does not preserve 2-cofiltered pseudolimits, in fact, it will preserve them only when \mathcal{C} is already a category of 2-pro-objects, in which case c is an equivalence.

2.2 Lemmas to compute with 2-pro-objects.

2.2.1 Definition.

1. Let $X \xrightarrow{f} Y$ be an arrow in $2\text{-Pro}(\mathcal{C})$. We say that a pair (r, φ) represents f , if φ is an invertible 2-cell $\pi_j f \xrightarrow{\varphi} r \pi_i$. That is, if we have the following diagram in $2\text{-Pro}(\mathcal{C})$:

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \pi_i \downarrow & \Downarrow \cong \varphi & \downarrow \pi_j \\ X_i & \xrightarrow{r} & Y_j \end{array}$$

2. Let $X \begin{array}{c} \xrightarrow{f} \\ \alpha \Downarrow \\ \xrightarrow{g} \end{array} Y$ and $X_i \begin{array}{c} \xrightarrow{r} \\ \theta \Downarrow \\ \xrightarrow{s} \end{array} Y_j$ be 2-cells in $2\text{-Pro}(\mathcal{C})$

and \mathcal{C} as in the following diagram:

$$\begin{array}{ccc}
 X & \xrightarrow{f} & Y \\
 \alpha \Downarrow & \xrightarrow{g} & \\
 \downarrow \pi_i & & \downarrow \pi_j \\
 X_i & \xrightarrow{r} & Y_j \\
 \theta \Downarrow & \xrightarrow{s} &
 \end{array}$$

We say that $(\theta, r, \varphi, s, \psi)$ represents α if (r, φ) represents f , (s, ψ) represents g , and the following diagram commutes in $2\text{-Pro}(\mathcal{C})$:

$$\begin{array}{ccc}
 \pi_j f & \xrightarrow[\cong]{\varphi} & r \pi_i \\
 \pi_j \alpha \Downarrow & & \Downarrow \theta \pi_i \\
 \pi_j g & \xrightarrow[\cong]{\psi} & s \pi_i
 \end{array}
 \quad i.e. \quad \theta \pi_i \circ \varphi = \psi \circ \pi_j \alpha$$

That is, $\theta \pi_i = \pi_j \alpha$ "modulo" a pair of invertible 2-cells φ, ψ .

Clearly, if α is invertible, then so is θ .

2.2.2 Proposition. Let $X = (X_i)_{i \in \mathcal{I}}$ and $Y = (Y_j)_{j \in \mathcal{J}}$ be any two objects in $2\text{-Pro}(\mathcal{C})$:

1. Let $X \xrightarrow{f} Y$, then, for any $j \in \mathcal{J}$ there is an $i \in \mathcal{I}$ and $X_i \xrightarrow{r} Y_j$ in \mathcal{C} , such that (r, id) represents f .

2. Let $X \xrightarrow[\alpha \Downarrow]{f} Y$, then, for any $j \in \mathcal{J}$ there is an $i \in \mathcal{I}$,

$X_i \xrightarrow[\theta \Downarrow]{r} Y_j$ in \mathcal{C} , and appropriate invertible 2-cells φ and ψ such that $(\theta, r, \varphi, s, \psi)$ represents α .

Proof. Consider $X \xrightarrow[\pi_j \alpha \Downarrow]{\pi_j f} Y_j$ and use formula 2.1.5 plus the constructions of pseudolimits and 2-filtered pseudocolimits, 1.2.7, 1.2.9. \square

2.2.3 Lemma. Let $X = (X_i)_{i \in \mathcal{I}} \in 2\text{-Pro}(\mathcal{C})$, let $X_i \xrightarrow{r} C$, $X_j \xrightarrow{s} C \in \mathcal{C}$, and let $X \xrightarrow[\alpha \Downarrow]{r\pi_i} C \in 2\text{-Pro}(\mathcal{C})$. Then, $\exists j \xrightarrow[u]{v} k$ and $X_k \xrightarrow[\theta \Downarrow]{rX_u} C$ such that:

$$\begin{array}{ccc}
 X & \xrightarrow{\pi_i} & X_i \\
 \pi_k \downarrow & \searrow \pi_u \Downarrow X_u & \downarrow r \\
 X_k & & C \\
 X_v \downarrow & \Downarrow \theta & \\
 X_j & \xrightarrow{s} & C
 \end{array}
 =
 \begin{array}{ccc}
 X & \xrightarrow{\pi_i} & X_i \\
 \pi_k \swarrow & & \downarrow \alpha \\
 X_k & \xleftarrow{\pi_v} & X_j \\
 X_v \searrow & & \downarrow r \\
 & & C \\
 & & \xrightarrow{s}
 \end{array}$$

$$\begin{array}{ccc}
 r\pi_i & \xrightarrow[r\pi_u]{\cong} & rX_u\pi_k \\
 \alpha \downarrow & & \downarrow \theta\pi_k \\
 s\pi_j & \xrightarrow[s\pi_v]{\cong} & sX_v\pi_k
 \end{array}
 \quad i.e. \quad \theta\pi_k \circ r\pi_u = s\pi_v \circ \alpha$$

Clearly, if α is invertible, then so is θ .

Proof. By formula 2.1.5 and the construction of 2-filtered pseudocolimits (1.2.9), α corresponds to a $(r, i) \xrightarrow{[u, \theta, v]} (s, j) \in \varinjlim_{i \in \mathcal{I}} \mathcal{C}(X_i, C)$. So,

$\exists j \xrightarrow[u]{v} k$ and $X_k \xrightarrow[\theta \Downarrow]{rX_u} C$ such that $\theta\pi_k \circ r\pi_u = s\pi_v \circ \alpha$, as we wanted to prove. \square

The following is an immediate consequence of [8, Lemma 2.2.]

2.2.4 Remark. If $i = j$, then one can choose $u = v$. \square

2.2.5 Lemma. Let $X = (X_i)_{i \in \mathcal{I}} \in 2\text{-Pro}(\mathcal{C})$ and $X_i \xrightarrow[\theta \Downarrow \theta' \Downarrow]{f} C \in \mathcal{C}$

be such that $\theta\pi_i = \theta'\pi_i$ in $2\text{-Pro}(\mathcal{C})$. Then $\exists i \xrightarrow{u} i'$ such that $\theta X_u = \theta' X_u$.

Proof. It follows from 2.1.5 and [8, Lemma 1.20.] \square

2.2.6 Lemma. Let $X \xrightarrow[\alpha \Downarrow \mathbf{g}]{\mathbf{f}} Y$ in $2\text{-Pro}(\mathcal{C})$ and $X_i \xrightarrow[\theta \Downarrow \theta' \Downarrow \mathbf{s}]{\mathbf{r}} Y_j$ in \mathcal{C} such that $(\theta, \mathbf{r}, \varphi, \mathbf{s}, \psi)$ and $(\theta', \mathbf{r}, \varphi, \mathbf{s}, \psi)$ both represents α . Then, there exists $i \xrightarrow{u} i' \in \mathcal{I}$ such that $\theta X_u = \theta' X_u$.

Proof. Since both $(\theta, \mathbf{r}, \varphi, \mathbf{s}, \psi)$ and $(\theta', \mathbf{r}, \varphi, \mathbf{s}, \psi)$ represents α , and φ, ψ are invertible, it follows that $\theta \pi_i = \theta' \pi_i$. Then, by 2.2.5, there exists $i \xrightarrow{u} i' \in \mathcal{I}$ such that $\theta X_u = \theta' X_u$. \square

2.2.7 Lemma. Let $X \xrightarrow[\alpha \Downarrow \mathbf{g}]{\mathbf{f}} Y \in 2\text{-Pro}(\mathcal{C})$, (\mathbf{r}, φ) representing \mathbf{f} , $X_i \xrightarrow{\mathbf{r}} Y_j$ and (\mathbf{s}, ψ) representing \mathbf{g} , $X_{i'} \xrightarrow{\mathbf{s}} Y_j$. Then, $\exists \begin{matrix} i \\ \xrightarrow{u} \\ i' \\ \xrightarrow{v} \end{matrix} k$ and $X_k \xrightarrow[\mathbf{s} X_v]{\mathbf{r} X_u} Y_j$ such that $(\theta, \mathbf{r} X_u, \mathbf{r} \pi_u \circ \varphi, \mathbf{s} X_v, \mathbf{s} \pi_v \circ \psi)$ represents α . Clearly, if α is invertible, then so is θ .

Proof. In lemma 2.2.3, take $C = Y_j$, and $\alpha = \psi \circ \pi_j \alpha \circ \varphi^{-1}$. Then,

$$\exists \begin{matrix} i \\ \xrightarrow{u} \\ i' \\ \xrightarrow{v} \end{matrix} k \text{ and } X_k \xrightarrow[\mathbf{s} X_v]{\mathbf{r} X_u} Y_j \text{ such that } \theta \pi_k \circ \mathbf{r} \pi_u \circ \varphi = \mathbf{s} \pi_v \circ \psi \circ \pi_j \alpha,$$

$$\begin{array}{ccccc} & & \mathbf{r} \pi_i & \xrightarrow{\mathbf{r} \pi_u} & \mathbf{r} X_u \pi_k & \xrightarrow{\theta \pi_k} & \mathbf{s} X_v \pi_k \\ & \nearrow \varphi & & & & & \\ \pi_j \mathbf{f} & & & & & & \\ & \searrow \pi_j \alpha & & & & & \\ & & \pi_j \mathbf{g} & \xrightarrow{\psi} & \mathbf{s} \pi_{i'} & & \end{array}$$

It is not difficult to check that $(\theta, \mathbf{r} X_u, \mathbf{r} \pi_u \circ \varphi, \mathbf{s} X_v, \mathbf{s} \pi_v \circ \psi)$ represents α . \square

From remark 2.2.4 we have:

2.2.8 Remark. If $i = i'$, then one can choose $u = v$. \square

2.3 2-cofiltered pseudolimits in $2\text{-Pro}(\mathcal{C})$.

Let \mathcal{J} be a small 2-filtered 2-category and $\mathcal{J}^{op} \xrightarrow{\mathbf{X}} 2\text{-Pro}(\mathcal{C})$ a 2-functor, $\mathbf{X}^j = (X_i^j)_{i \in \mathcal{I}_j}$, $\mathcal{I}_j^{op} \xrightarrow{\mathbf{X}^j} \mathcal{C}$. Recall (2.1.8) that for each

j in \mathcal{J} , \mathbf{X}^j is equipped with a pseudolimit pseudocone $\{\pi_i^j\}_{i \in \mathcal{I}_j}$, $\{\pi_u^j\}_{i \xrightarrow{u} i' \in \mathcal{I}_j}$ for the 2-functor \mathbf{X}^j .

We are going to construct a 2-pro-object which is going to be the pseudolimit of \mathbf{X} in $2\text{-Pro}(\mathcal{C})$. First we construct its index category

2.3.1 Definition. Let $\mathcal{K}_{\mathbf{X}}$ be the 2-category consisting on:

1. 0-cells of $\mathcal{K}_{\mathbf{X}}$: (i, j) , where $j \in \mathcal{J}$, $i \in \mathcal{I}_j$.
2. 1-cells of $\mathcal{K}_{\mathbf{X}}$: $(i, j) \xrightarrow{(a, r, \varphi)} (i', j')$, where $j \xrightarrow{a} j' \in \mathcal{J}$, $\mathbf{X}_{i'}^{j'} \xrightarrow{r} \mathbf{X}_i^j$ are such that (r, φ) represents \mathbf{X}^a .
3. 2-cells of $\mathcal{K}_{\mathbf{X}}$: $(a, r, \varphi) \xrightarrow{(\alpha, \theta)} (b, s, \psi)$, where $a \xrightarrow{\alpha} b \in \mathcal{J}$ and $(\theta, r, \varphi, s, \psi)$ represents \mathbf{X}^α .

The 2-category structure is given as follows:

$$\begin{array}{ccccc}
 & \xrightarrow{(a, r, \varphi)} & & \xrightarrow{(a', r', \varphi')} & \\
 & \Downarrow(\alpha, \theta) & & \Downarrow(\alpha', \theta') & \\
 (i, j) & \xrightarrow{(b, s, \psi)} & (i', j') & \xrightarrow{(b', s', \psi')} & (i'', j'') \\
 & \Downarrow(\beta, \eta) & & \Downarrow(\beta', \eta') & \\
 & \xrightarrow{(c, t, \phi)} & & \xrightarrow{(c', t', \phi')} &
 \end{array}$$

1. $(a', r', \varphi')(a, r, \varphi) = (a'a, rr', r\varphi' \circ \varphi \mathbf{X}^{a'})$
2. $(\alpha', \theta')(\alpha, \theta) = (\alpha'\alpha, \theta\theta')$
3. $(\beta, \eta) \circ (\alpha, \theta) = (\beta \circ \alpha, \eta \circ \theta)$

One can easily check that the structure so defined is indeed a 2-category, which is clearly small.

2.3.2 Proposition. The 2-category $\mathcal{K}_{\mathbf{X}}$ is 2-filtered.

Proof. F0. Let $(i, j), (i', j') \in \mathcal{K}_{\mathbf{X}}$. Since \mathcal{J} is 2-filtered, $\exists \begin{array}{c} j \\ \xrightarrow{a} \\ j' \xrightarrow{b} \end{array} j''$.

By 2.2.2, $\exists \mathbf{X}_{i_1}^{j''} \xrightarrow{r_1} \mathbf{X}_i^j$ and $\mathbf{X}_{i_2}^{j''} \xrightarrow{r_2} \mathbf{X}_{i'}^{j'}$ such that (r_1, id) represents

X^a and (r_2, id) represents X^b . Since $\mathcal{I}_{j''}$ is 2-filtered, $\exists \begin{matrix} i_1 \\ \xrightarrow{u} \\ i_2 \end{matrix} \xrightarrow{v} i''$. Then, we have the following situation in \mathcal{K}_X which proves $F0$.:

$$\begin{array}{ccc} (i, j) & \xrightarrow{(a, r_1 X_u^{j''}, r_1 \pi_u^{j''})} & (i'', j'') \\ & \searrow & \nearrow \\ (i', j') & \xrightarrow{(b, r_2 X_v^{j''}, r_2 \pi_v^{j''})} & \end{array}$$

F1. Let $(i, j) \xrightarrow[(b, s, \psi)]{(a, r, \varphi)} (i', j') \in \mathcal{K}_X$. Since \mathcal{J} is 2-filtered,

$\exists j' \xrightarrow{c} j''$ and an invertible 2-cell $ca \xrightarrow{\alpha} cb$. By 2.2.2, $\exists X_k^{j''} \xrightarrow{t} X_{i'}^{j''}$ such that (t, id) represents X^c . Then $(rt, \varphi X^c)$ represents X^{ca} and $(st, \psi X^c)$ represents X^{cb} , so, by 2.2.7, there exists $k \xrightarrow{w} i'' \in \mathcal{I}_{j''}$ and an invertible 2-cell $rtX_w^{j''} \xrightarrow{\theta} stX_w^{j''}$ such that $(\theta, rtX_w^{j''}, rt\pi_w \circ \varphi X^c, stX_w^{j''}, st\pi_w \circ \psi X^c)$ represents X^α . Then

$$\text{we have an invertible 2-cell in } \mathcal{K}_X \quad (i, j) \xrightarrow[(c, tX_w^{j''}, t\pi_w)(b, s, \psi)]{(c, tX_w^{j''}, t\pi_w)(a, r, \varphi)} \downarrow (\alpha, \theta) \xrightarrow{(i'', j'')} \downarrow (c, tX_w^{j''}, t\pi_w)$$

which proves $F1$.

F2. Let $(i, j) \xrightarrow[(b, s, \psi)]{(a, r, \varphi)} \downarrow (\alpha, \theta) \downarrow (\alpha', \theta') \xrightarrow{(i', j')} \downarrow (b, s, \psi)$. Since \mathcal{J} is

2-filtered, $\exists j' \xrightarrow{c} j'' \in \mathcal{J}$ such that $c\alpha = c\alpha'$. Also, by 2.2.2, $\exists X_k^{j''} \xrightarrow{t} X_{i'}^{j''}$ such that (t, id) represents X^c . Then, it is easy to check that (t, t, id, t, id) represents X^c and therefore we have that $(\theta t, rt, \varphi X^c, st, \psi X^c)$ and $(\theta' t, rt, \varphi X^c, st, \psi X^c)$ both represent $X^{c\alpha}$. Then, by 2.2.6, $\exists k \xrightarrow{w} i'' \in \mathcal{I}_{j''}$ such that $\theta tX_w^{j''} = \theta' tX_w^{j''}$, so $(c, tX_w^{j''}, t\pi_w)(\alpha, \theta) = (c, tX_w^{j''}, t\pi_w)(\alpha', \theta')$, which proves $F2$. \square

2.3.3 Theorem. Let \tilde{X} be the 2-pro-object $\mathcal{K}_X^{op} \xrightarrow{\tilde{X}} \mathcal{C}$ defined by $\tilde{X}_{(i, j)} = X_i^j$, $\tilde{X}_{(a, r, \varphi)} = r$, and $\tilde{X}_{(\alpha, \theta)} = \theta$. Then the following equation holds in 2-Pro(\mathcal{C}):

$$\tilde{X} = \underset{j \in \mathcal{J}^{op}}{\text{L}\lim} X^j$$

Proof. Let $Z \in 2\text{-Pro}(\mathcal{C})$, and $\{Z \xrightarrow{h_j} X^j\}_{j \in \mathcal{J}}$, $\{h_j \xrightarrow{h_a} X^a h_{j'}\}_{j \xrightarrow{a} j' \in \mathcal{J}}$ be a pseudocone for X with vertex Z (1.2.1). Given $(i, j) \xrightarrow{(a, r, \varphi)} (i', j') \in \mathcal{K}_X$, check that the definitions $h_{(i, j)} = \pi_i^j h_j$ and $h_{(a, r, \varphi)} = \varphi h_{j'} \circ \pi_i^j h_a$ determine a pseudocone for $c\tilde{X}$ with vertex Z . It is straightforward to check that this extends to a functor, that we denote p (for the isomorphism below see 2.1.9):

$$\text{PC}_{2\text{-Pro}(\mathcal{C})}(Z, X) \xrightarrow{p} \text{PC}_{2\text{-Pro}(\mathcal{C})}(Z, c\tilde{X}) \cong 2\text{-Pro}(\mathcal{C})(Z, \tilde{X})$$

The theorem follows if p is an isomorphism. In the sequel we prove that, in fact, p is an isomorphism. Let $Z \in 2\text{-Pro}(\mathcal{C})$, and

$$\{h_{(i, j)} \xrightarrow{h_{(a, r, \varphi)}} \tilde{X}_{(a, r, \varphi)} h_{(i', j')} = r h_{(i, j)}\}_{(i, j) \xrightarrow{(a, r, \varphi)} (i', j') \in \mathcal{K}_X}, \quad \{Z \xrightarrow{h_{(i, j)}} X_i^j\}_{(i, j) \in \mathcal{K}_X}$$

be a pseudocone for $c\tilde{X}$ with vertex Z (1.2.1).

1. p is bijective on objects:

Check that for each $j \in \mathcal{J}$, $\{Z \xrightarrow{h_{(i, j)}} X_i^j\}_{i \in \mathcal{I}_j}$ together with $\{h_u = h_{(j, X_u^j, \pi_u^j)} : h_{(i, j)} \implies X_u^j h_{(i', j')}\}_{i \xrightarrow{u} i' \in \mathcal{I}_j}$ is a pseudocone for X^j . Then, since $X^j \xrightarrow{\pi_i^j} X_i^j$ is a pseudolimit pseudocone, it follows that there exists a unique $Z \xrightarrow{h_j} X^j$ such that

$$(2.3.4) \quad \forall i \in \mathcal{I}_j \quad \pi_i^j h_j = h_{(i, j)} \quad \text{and} \quad \forall i \xrightarrow{u} i' \in \mathcal{I}_j \quad \pi_u^j h_j = h_u.$$

It only remains to define the 2-cells of the pseudocone structure. That is, for each $j \xrightarrow{a} j' \in \mathcal{J}$, we need invertible 2-cells $h_j \xrightarrow{h_a} h_{j'} X^a$, such that $\{h_j\}_{j \in \mathcal{J}}$ together with $\{h_a\}_{j \xrightarrow{a} j' \in \mathcal{J}}$ form a pseudocone for X with vertex Z .

Consider the pseudocone $\{X^j \xrightarrow{\pi_i^j} X_i^j\}_{i \in \mathcal{I}_j}$. Then the composites $\pi_i^j h_j$, $\pi_i^j X^a h_{j'}$, determine two pseudocones $\{Z \xrightarrow[\pi_i^j X^a h_{j'}]{\pi_i^j h_j} X_i^j\}_{i \in \mathcal{I}_j}$ for X^j with vertex Z .

Claim 1 *Let (r, φ) and (s, ψ) be two pairs representing X^a as follows:*

$$\begin{array}{ccc} X^{j'} & \xrightarrow{X^a} & X^j \\ \pi_{i'}^{j'} \downarrow & \Downarrow \cong \varphi & \downarrow \pi_i^j \\ X_{i'}^{j'} & \xrightarrow{r} & X_i^j \end{array} \quad \begin{array}{ccc} X^{j'} & \xrightarrow{X^a} & X^j \\ \pi_{i''}^{j'} \downarrow & \Downarrow \cong \psi & \downarrow \pi_i^j \\ X_{i''}^{j'} & \xrightarrow{s} & X_i^j \end{array}$$

Then, $\varphi^{-1}h_{j'} \circ h_{(a,r,\varphi)} = \psi^{-1}h_{j'} \circ h_{(a,s,\psi)}$ (proof below).

Claim 2 *For each $i \in \mathcal{I}_j$, let (r, φ) be a pair representing X^a , and set $\rho_i = \varphi^{-1}h_{j'} \circ h_{(a,r,\varphi)}$. Then, $\{\rho_i\}_{i \in \mathcal{I}_j}$ determines an isomorphism of*

pseudocones $\{Z \xrightarrow[\pi_i^j X^a h_{j'}]{\rho_i \Downarrow} X_i^j\}_{i \in \mathcal{I}_j}$ (proof below).

Since $X^j \xrightarrow{\pi_i^j} X_i^j$ is a pseudolimit pseudocone, the functor $2\text{-Pro}(\mathcal{C})(Z, X^j) \xrightarrow{(\pi^j)^*} \text{PC}_{2\text{-Pro}(\mathcal{C})}(Z, X^j)$ is an isomorphism of categories. Then, from Claim 2 it follows that there are invertible 2-cells $Z \xrightarrow[\pi_i^j X^a h_{j'}]{h_j} X^j \in 2\text{-Pro}(\mathcal{C})$ such that $\rho_i = \pi_i^j h_a \forall i \in \mathcal{I}_j$. It can be

checked that in fact $\{Z \xrightarrow{h_j} X^j\}_{j \in \mathcal{J}}$ with $\{h_j \xrightarrow{h_a} h_{j'} X^a\}_{j \rightarrow j' \in \mathcal{J}}$ is a pseudocone over X .

2. *p is full and faithful:*

Let $\{Z \xrightarrow[\pi_i^j m_j]{\rho_{(i,j)} \Downarrow} X_i^j\}_{(i,j) \in \mathcal{K}_X}$ be a morphism of pseudocones for \tilde{X} . It

is easy to check that for each $j \in \mathcal{J}$, $\{Z \xrightarrow[\pi_i^j m_j]{\rho_{(i,j)} \Downarrow} X_i^j\}_{i \in \mathcal{I}_j}$ is a morphism

of pseudocones for X^j . Then arguing as above, there exists a unique morphism $Z \xrightarrow[\pi_i^j m_j]{h_j} X^j \in 2\text{-Pro}(\mathcal{C})$ such that $\forall i \in \mathcal{I}_j$, $\pi_i^j \rho_j = \rho_{(i,j)}$.

It can be checked that $\{\rho_j\}_{j \in \mathcal{J}}$ is a morphism of pseudocones. This

proves the assertion. \square

Proof of Claim 1. First assume that $i' = i''$ and (r, φ) , (s, ψ) are related by a 2-cell $(i, j) \begin{array}{c} \xrightarrow{(a,r,\varphi)} \\ \xrightarrow{(a,\theta)\Downarrow} \\ \xrightarrow{(a,s,\psi)} \end{array} (i', j')$ in \mathcal{K}_X . Then:

$$(\psi^{-1}h_{j'}) \circ (h_{(a,s,\psi)}) = (\psi^{-1}h_{j'}) \circ (\theta h_{(i',j')}) \circ h_{(a,r,\varphi)} = (\varphi^{-1}h_{j'}) \circ (h_{(a,r,\varphi)}),$$

the first equality by the pseudocone axiom PC2 (Definition 1.2.1), and the second because θ represents *id* (the identity of X^a).

The general case reduces to this one as follows:

We have $(i, j) \begin{array}{c} \xrightarrow{(a,r,\varphi)} \\ \xrightarrow{(a,s,\psi)} \end{array} \begin{array}{c} (i', j') \\ (i'', j') \end{array}$. Take $\begin{array}{c} i' \\ i'' \end{array} \begin{array}{c} \xrightarrow{u} \\ \xrightarrow{v} \end{array} k$ in \mathcal{I}_j . This yields a particular instance of lemma 2.2.7:

$$\begin{array}{ccc} X^{j'} & \xrightarrow{X^a} & X^j \\ \downarrow \pi_k & \begin{array}{c} \xrightarrow{id \Downarrow} \\ \xrightarrow{X^a} \end{array} & \downarrow \pi_i \\ X_k^{j'} & \xrightarrow{\begin{array}{c} rX_u^{j'} \\ sX_v^{j'} \end{array}} & X_i^j \end{array}$$

with $(rX_u^{j'}, (r\pi_u^{j'}) \circ \varphi)$ and $(sX_v^{j'}, (s\pi_v^{j'}) \circ \psi)$ both representing X^a .

It follows there exists $k \xrightarrow{w} k'$ and $X_{k'}^{j'} \begin{array}{c} \xrightarrow{rX_u^{j'} X_w^{j'}} \\ \xrightarrow{\theta \Downarrow} \\ \xrightarrow{sX_v^{j'} X_w^{j'}} \end{array} X_i^j$ such that

$(\theta, rX_u^{j'} X_w^{j'}, rX_u^{j'} \pi_w^{j'} \circ r\pi_u^{j'} \circ \varphi, sX_v^{j'} X_w^{j'}, sX_v^{j'} \pi_w^{j'} \circ s\pi_v^{j'} \circ \psi)$ represents *id* (the identity of X^a).

Considering $(rX_u^{j'} X_w^{j'}, rX_u^{j'} \pi_w^{j'} \circ r\pi_u^{j'} \circ \varphi)$ and $(sX_v^{j'} X_w^{j'}, sX_v^{j'} \pi_w^{j'} \circ s\pi_v^{j'} \circ \psi)$ both representing X^a , we have a situation that corresponds to the previous case. Thus:

$$\begin{aligned} & (\varphi^{-1}h_{j'}^{-1} \circ r(\pi_u^{j'})^{-1} \circ rX_u^{j'} (\pi_w^{j'})^{-1}) h_{j'} \circ rh_{(j', X_u^{j'} X_w^{j'}, X_u^{j'} \pi_w^{j'})} \circ h_{(a,r,\varphi)} = \\ & = (\psi^{-1}h_{j'}^{-1} \circ s(\pi_v^{j'})^{-1} \circ sX_v^{j'} (\pi_w^{j'})^{-1}) h_{j'} \circ sh_{(j', X_v^{j'} X_w^{j'}, X_v^{j'} \pi_w^{j'})} \circ h_{(a,s,\psi)}. \end{aligned}$$

From 2.3.4, it follows that $(r(\pi_u^{j'})^{-1} \circ r\mathbf{X}_u^{j'}(\pi_w^{j'})^{-1})\mathbf{h}_{j'} \circ \mathbf{r}\mathbf{h}_{(j', \mathbf{X}_u^{j'}, \mathbf{X}_w^{j'}, \mathbf{X}_u^{j'} \pi_w^{j'})}$ and $(s(\pi_v^{j'})^{-1} \circ s\mathbf{X}_v^{j'}(\pi_w^{j'})^{-1})\mathbf{h}_{j'} \circ \mathbf{s}\mathbf{h}_{(j', \mathbf{X}_v^{j'}, \mathbf{X}_w^{j'}, \mathbf{X}_v^{j'} \pi_w^{j'})}$ are identities. So $\varphi^{-1}\mathbf{h}_{j'} \circ \mathbf{h}_{(a,r,\varphi)} = \psi^{-1}\mathbf{h}_{j'} \circ \mathbf{h}_{(a,s,\psi)}$ as we wanted to prove. \square

Proof of Claim 2. Given any $i \xrightarrow{u} k \in \mathcal{I}_j$, we have to check the PCM equation in 1.2.1. Given the pair (s, ψ) used to define ρ_k , it is possible to choose a pair (r, φ) to define ρ_i in such a way that the equation holds. This arguing is justified by Claim 1. \square

2.3.5 Corollary. *2-Pro(\mathcal{C}) is closed under small 2-cofiltered pseudolimits. Considering the equivalence in 2.1.2, it follows that the inclusion $\mathcal{H}\text{om}(\mathcal{C}, \text{Cat})_{fc} \subset \mathcal{H}\text{om}(\mathcal{C}, \text{Cat})$ is closed under small 2-filtered pseudocolimits* \square

2.4 Universal property of 2-Pro(\mathcal{C})

In this subsection we prove for 2-pro-objects the universal property established for pro-objects in [1, Ex. I, Prop. 8.7.3]. Consider the 2-functor $\mathcal{C} \xrightarrow{c} 2\text{-Pro}(\mathcal{C})$ of Corollary 2.1.7 and a 2-pro-object $\mathbf{X} = (\mathbf{X}_i)_{i \in \mathcal{I}}$. Given a 2-functor $\mathcal{C} \xrightarrow{F} \mathcal{E}$ into a 2-category closed under small 2-cofiltered pseudolimits, we can naively extend F into a 2-cofiltered pseudolimit preserving 2-functor $2\text{-Pro}(\mathcal{C}) \xrightarrow{\widehat{F}} \mathcal{E}$ by defining $\widehat{F}\mathbf{X} = \varprojlim_{i \in \mathcal{I}} F\mathbf{X}_i$. This is just part of a 2-equivalence of 2-categories that we develop with the necessary precision in this subsection. First the universal property should be wholly established for $\mathcal{E} = \text{Cat}$, and only afterwards can be lifted to any 2-category \mathcal{E} closed under small 2-cofiltered pseudolimits.

2.4.1 Lemma. *Let \mathcal{C} be a 2-category and $F : \mathcal{C} \rightarrow \text{Cat}$ a 2-functor. Then, there exists a 2-functor $\widehat{F} : 2\text{-Pro}(\mathcal{C}) \rightarrow \text{Cat}$ that preserves small 2-cofiltered pseudolimits, and an isomorphism $\widehat{F}c \xrightarrow{\cong} F$ in $\mathcal{H}\text{om}(\mathcal{C}, \text{Cat})$.*

Proof. Let $\mathbf{X} = (\mathbf{X}_i)_{i \in \mathcal{I}} \in 2\text{-Pro}(\mathcal{C})$ be a 2-pro-object. Define:

$$\begin{aligned} \widehat{FX} &= (\mathcal{H}om(\mathcal{C}, \mathcal{C}at)(-, F) \circ L)X = \mathcal{H}om(\mathcal{C}, \mathcal{C}at)(\varinjlim_{i \in \mathcal{I}} \mathcal{C}(X_i, -), F) \xrightarrow{\cong} \\ &\xrightarrow{\cong} \varprojlim_{i \in \mathcal{I}} \mathcal{H}om(\mathcal{C}, \mathcal{C}at)(\mathcal{C}(X_i, -), F) \xrightarrow{\cong} \varprojlim_{i \in \mathcal{I}} FX_i. \end{aligned}$$

Where L is the 2-functor of 2.1.2, the first isomorphism is by definition of pseudocolimit 1.2.3, and the second is the Yoneda isomorphism 1.1.16. Since it is a 2-equivalence, the 2-functor L preserves any pseudocolimit. Then by Corollary 2.3.5 it follows that the composite $\mathcal{H}om(\mathcal{C}, \mathcal{C}at)(-, F) \circ L$ preserves small 2-cofiltered pseudocolimits \square

2.4.2 Theorem. *Let \mathcal{C} be any 2-category. Then, pre-composition with $\mathcal{C} \xrightarrow{c} 2\text{-Pro}(\mathcal{C})$ is a 2-equivalence of 2-categories:*

$$\mathcal{H}om(2\text{-Pro}(\mathcal{C}), \mathcal{C}at)_+ \xrightarrow{c^*} \mathcal{H}om(\mathcal{C}, \mathcal{C}at)$$

(where $\mathcal{H}om(2\text{-Pro}(\mathcal{C}), \mathcal{C}at)_+$ stands for the full subcategory whose objects are those 2-functors that preserve small 2-cofiltered pseudocolimits).

Proof. We check that the 2-functor c^* is essentially surjective on objects and 2-fully-faithful:

Essentially surjective on objects: It follows from lemma 2.4.1.

2-fully-faithful: We check that if F and G are 2-functors from $2\text{-Pro}(\mathcal{C})$ to $\mathcal{C}at$ that preserve small 2-cofiltered pseudocolimits, then

$$(2.4.3) \quad \mathcal{H}om(2\text{-Pro}(\mathcal{C}), \mathcal{C}at)_+(F, G) \xrightarrow{c^*} \mathcal{H}om(\mathcal{C}, \mathcal{C}at)(Fc, Gc)$$

is an isomorphism of categories.

Let $Fc \xrightarrow[\eta_c]{\theta_c} Gc \in \mathcal{H}om(\mathcal{C}, \mathcal{C}at)(Fc, Gc)$. It can be easily checked

that the composites $\{FX \xrightarrow{F\pi_i} FX_i \xrightarrow[\eta_{X_i}]{\theta_{X_i}} GX_i\}_{i \in \mathcal{I}}$ determine two pseudocones for GX together with a morphism of pseudocones. Since G preserves small 2-cofiltered pseudocolimits, post-composing with $GX \xrightarrow{G\pi_i} GX_i$ is an isomorphism of categories $\mathcal{C}at(FX, GX) \xrightarrow{(G\pi)^*} \mathcal{P}C_{\mathcal{C}at}(FX, GX)$. It follows there exists a unique 2-cell in $\mathcal{C}at$, $FX \xrightarrow[\eta'_X]{\theta'_X} GX$, such that

$G\pi_i\theta'_X = \theta_{X_i}F\pi_i$, $G\pi_i\eta'_X = \eta_{X_i}F\pi_i$, and $G\pi_i\mu'_X = \mu_{X_i}F\pi_i$, $\forall i \in \mathcal{I}$. It is not difficult to check that θ'_X , η'_X are in fact 2-natural on X , and that μ'_X is a modification. Clearly $\theta'c = \theta$, $\eta'c = \eta$, and $\mu'c = \mu$. Thus 2.4.3 is an isomorphism of categories. \square

2.4.4 Lemma. *Let \mathcal{C} be a 2-category, \mathcal{E} a 2-category closed under small 2-cofiltered pseudolimits and $F : \mathcal{C} \rightarrow \mathcal{E}$ a 2-functor. Then, there exists a 2-functor $\widehat{F} : 2\text{-Pro}(\mathcal{C}) \rightarrow \mathcal{E}$ that preserves small 2-cofiltered pseudolimits, and an isomorphism $\widehat{F}c \xrightarrow{\cong} F$ in $\mathcal{H}om(\mathcal{C}, \mathcal{E})$.*

Proof. If $X = (X_i)_{i \in \mathcal{I}} \in 2\text{-Pro}(\mathcal{C})$, define $\widehat{F}X = \varprojlim_{i \in \mathcal{I}^{op}} FX_i$. We will prove that this is the object function part of a 2-functor, and that this 2-functor has the rest of the properties asserted in the proposition.

Consider the composition $y_{(-)}F : \mathcal{C} \xrightarrow{F} \mathcal{E} \xrightarrow{y_{(-)}} \mathcal{H}om(\mathcal{E}^{op}, \mathcal{C}at)$, where $y_{(-)}$ is the Yoneda 2-functor (1.1.15). Under the isomorphism 1.1.14 this corresponds to a 2-functor $\widehat{(-)} : \mathcal{C} \rightarrow \mathcal{H}om(\mathcal{C}, \mathcal{C}at)$. Composing this 2-functor with a quasi-inverse $\widetilde{(-)}$ for the 2-equivalence in 2.4.2, we obtain a 2-functor $\mathcal{E}^{op} \rightarrow \mathcal{H}om(2\text{-Pro}(\mathcal{C}), \mathcal{C}at)_+$, which in turn corresponds to a 2-functor $2\text{-Pro}(\mathcal{C}) \xrightarrow{\widetilde{F}} \mathcal{H}om(\mathcal{E}^{op}, \mathcal{C}at)$. The 2-functor \widetilde{F} preserves small 2-cofiltered pseudolimits because they are computed pointwise in $\mathcal{H}om(\mathcal{E}^{op}, \mathcal{C}at)$ (1.2.4). Chasing the isomorphisms shows that we have the following diagram:

$$(2.4.5) \quad \begin{array}{ccc} \mathcal{C} & \xrightarrow{c} & 2\text{-Pro}(\mathcal{C}) \\ \downarrow F & \Downarrow \cong & \downarrow \widetilde{F} \\ \mathcal{E} & \xrightarrow{y_{(-)}} & \mathcal{H}om(\mathcal{E}^{op}, \mathcal{C}at) \end{array}$$

$\widetilde{F}c \xrightarrow{\cong} y_{(-)}F$

Consider the following chain of isomorphisms (the first and the third because \widetilde{F} and $y_{(-)}$ preserve pseudolimits (1.2.6), and the middle one given by 2.4.5):

$$\widetilde{F}X = \widetilde{F}\varprojlim_{i \in \mathcal{I}} X_i \xrightarrow{\cong} \varprojlim_{i \in \mathcal{I}} \widetilde{F}cX_i \xrightarrow{\cong} \varprojlim_{i \in \mathcal{I}} y_{(-)}FX_i \xleftarrow{\cong} y_{(-)}\varprojlim_{i \in \mathcal{I}} FX_i.$$

This shows that $\widetilde{F}X$ is in the essential image of $y_{(-)}$. Since $y_{(-)}$ is 2-fully faithful (1.1.17), it follows there is a factorization $y_{(-)}\widehat{F} \xrightarrow{\cong} \widetilde{F}$,

given by a 2-functor $2\text{-Pro}(\mathcal{C}) \xrightarrow{\widehat{F}} \mathbf{E}$. Clearly \widehat{F} preserves small 2-cofiltered pseudolimits. We have $y_{(-)}\widehat{F}c \xrightarrow{\cong} \widetilde{F}c \xrightarrow{\cong} y_{(-)}\widehat{F}$. Finally, the fully faithfulness of $y_{(-)}$ provides an isomorphism $y_{(-)}\widehat{F} \xrightarrow{\cong} F$. This finishes the proof. \square

Exactly the same proof of theorem 2.4.2 applies with an arbitrary 2-category \mathcal{E} in place of $\mathcal{C}at$, and we have:

2.4.6 Theorem. *Let \mathcal{C} be any 2-category, and \mathcal{E} a 2-category closed under small 2-cofiltered pseudolimits. Then, pre-composition with $\mathcal{C} \xrightarrow{c} 2\text{-Pro}(\mathcal{C})$ is a 2-equivalence of 2-categories:*

$$\mathcal{H}om(2\text{-Pro}(\mathcal{C}), \mathcal{E})_+ \xrightarrow{c^*} \mathcal{H}om(\mathcal{C}, \mathcal{E})$$

Where $\mathcal{H}om(2\text{-Pro}(\mathcal{C}), \mathcal{E})_+$ stands for the full subcategory whose objects are those 2-functors that preserve small 2-cofiltered pseudolimits. \square

From theorem 2.4.6 it follows automatically the pseudo-functoriality of the assignment of the 2-category $2\text{-Pro}(\mathcal{C})$ to each 2-category \mathcal{C} , and in such a way that c becomes a pseudonatural transformation. But we can do better:

If we put $\mathcal{E} = 2\text{-Pro}(\mathcal{D})$ in 2.4.6 it follows there is a 2-functor (post-composing with c followed by a quasi-inverse in 2.4.6)

$$(2.4.7) \quad \mathcal{H}om(\mathcal{C}, \mathcal{D}) \xrightarrow{\widehat{(-)}} \mathcal{H}om(2\text{-Pro}(\mathcal{C}), 2\text{-Pro}(\mathcal{D}))_+,$$

and for each 2-functor $\mathcal{C} \xrightarrow{F} \mathcal{D}$, a diagram:

$$(2.4.8) \quad \begin{array}{ccc} 2\text{-Pro}(\mathcal{C}) & \xrightarrow{\widehat{F}} & 2\text{-Pro}(\mathcal{D}) \\ \uparrow c & \Downarrow \cong & \uparrow c \\ \mathcal{C} & \xrightarrow{F} & \mathcal{D} \end{array}$$

Given any 2-pro-object $X \in 2\text{-Pro}(\mathcal{C})$, set $2\text{-Pro}(F)(X) = \widehat{F}X$. It is straightforward to check that this determines a 2-functor

$$2\text{-Pro}(\mathcal{C}) \xrightarrow{2\text{-Pro}(F)} 2\text{-Pro}(\mathcal{D})$$

making diagram 2.4.8 commutative. It follows we have an isomorphism $\widehat{F}X \xrightarrow{\cong} 2\text{-Pro}(\mathbf{F})(X)$ 2-natural in X . This shows that the 2-functor $2\text{-Pro}(\mathbf{F})$ preserves small 2-cofiltered pseudolimits because \widehat{F} does. Also, it follows that $2\text{-Pro}(\mathbf{F})$ determines a 2-functor as in 2.4.7. In conclusion, denoting now by 2-CAT the 2-category of locally small 2-categories (see 1.1.10) we have:

2.4.9 Theorem. *The definition $2\text{-Pro}(\mathbf{F})(X) = \widehat{F}X$ determines a 2-functor*

$$2\text{-Pro}(-) : 2\text{-CAT} \longrightarrow 2\text{-CAT}_+$$

in such a way that c becomes a 2-natural transformation (where 2-CAT_+ is the full sub 2-category of locally small 2-categories closed under small 2-cofiltered pseudo limits and small pseudolimit preserving 2-functors). \square

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